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Hidden Markov Models

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Nov 29th, 2001

A Markov System

 $\left(\mathbf{s}_{2}\right)$

Has N states, called s_1 , s_2 ... s_N There are discrete timesteps, t=0, t=1, ...



 $\left(\mathbf{s}_{3}\right)$

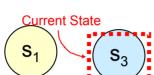
N = 3

t=0

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A Markov System





N = 3

t=0

 $q_t = q_0 = s_3$

Has N states, called s_1 , s_2 .. s_N

There are discrete timesteps, t=0, t=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it q_t

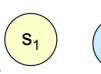
Note: $q_t \in \{s_1, s_2 ... s_N\}$

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Hidden Markov Models: Slide 3

A Markov System

Current State S₂



N = 3

t=1

 $q_t = q_1 = s_2$

Has N states, called s_1 , s_2 .. s_N

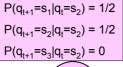
There are discrete timesteps, t=0, t=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it q_t

Note: $q_t \in \{s_1, s_2 ... s_N \}$

Between each timestep, the next state is chosen randomly.

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$$P(q_{t+1}=s_1|q_t=s_1) = 0$$

$$P(q_{t+1}=s_2|q_t=s_1) = 0$$

$$P(q_{t+1}=s_3|q_t=s_1) = 1$$

$$\eta_t = s_1 = 1$$

$$q_t = q_1 = s_2$$



 S_3

 $P(q_{t+1}=s_2|q_t=s_3) = 2/3$

 $P(q_{t+1}=s_3|q_t=s_3)=0$

A Markov System

Has N states, called s_1 , s_2 .. s_N

There are discrete timesteps. t=0, t=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it q,

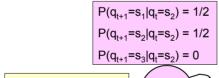
Note:
$$q_t \in \{s_1, s_2 ... s_N\}$$

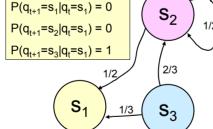
Between each timestep, the next $P(q_{t+1}=s_1|q_t=s_3) = 1/3$ | state is chosen randomly.

> The current state determines the probability distribution for the next state.

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Hidden Markov Models: Slide 5





N = 3

$$q_t = q_1 = s_2$$

 $P(q_{t+1}=s_1|q_t=s_3) = 1/3$

$$P(q_{t+1}=s_2|q_t=s_3) = 2/3$$

$$P(q_{t+1}=s_3|q_t=s_3)=0$$

Often notated with arcs between states

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A Markov System

Has N states, called s_1 , s_2 .. s_N

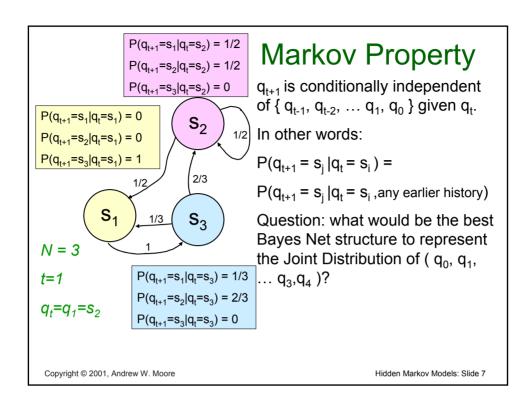
There are discrete timesteps, 1/2 t=0, t=1, ...

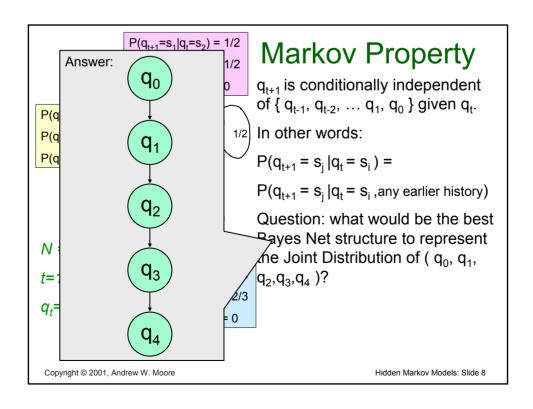
> On the t'th timestep the system is in exactly one of the available states. Call it q_t

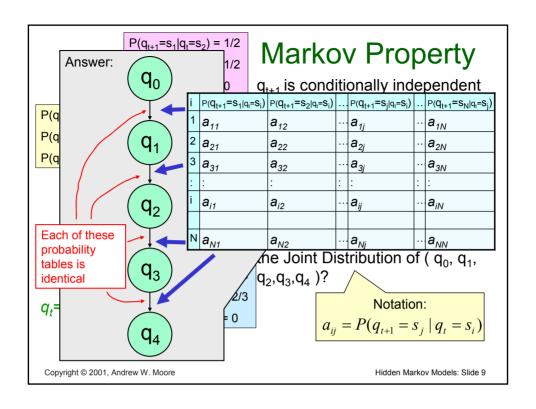
Note:
$$q_t \in \{s_1, s_2 ... s_N\}$$

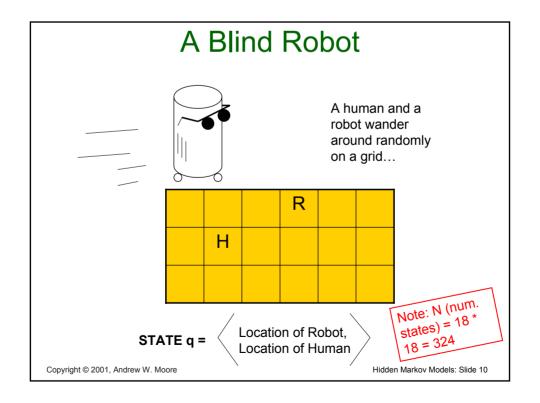
Between each timestep, the next state is chosen randomly.

The current state determines the probability distribution for the next state.

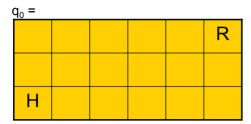








Dynamics of System

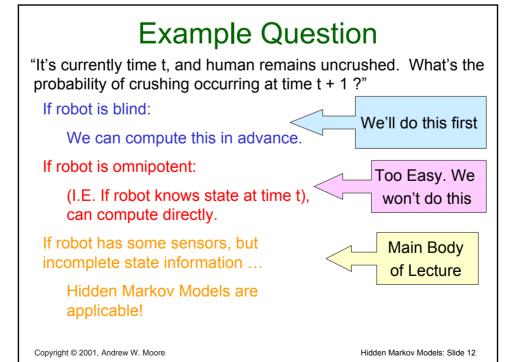


Each timestep the human moves randomly to an adjacent cell. And Robot also moves randomly to an adjacent cell.

Typical Questions:

- "What's the expected time until the human is crushed like a bug?"
- "What's the probability that the robot will hit the left wall before it hits the human?"
- "What's the probability Robot crushes human on next time step?"

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What is $P(q_t = s)$? slow, stupid answer

Step 1: Work out how to compute P(Q) for any path Q $= q_1 q_2 q_3 ... q_t$

Given we know the start state q₁

$$\begin{split} P(q_1 \; q_2 \; ... \; q_t) &= P(q_1 \; q_2 \; ... \; q_{t-1}) \; P(q_t | q_1 \; q_2 \; ... \; q_{t-1}) \\ &= P(q_1 \; q_2 \; ... \; q_{t-1}) \; P(q_t | q_{t-1}) & \textit{WHY?} \\ &= P(q_2 | q_1) P(q_3 | q_2) ... P(q_t | q_{t-1}) \end{split}$$

Step 2: Use this knowledge to get $P(q_t = s)$

$$P(q_t = s) = \sum_{Q \in Paths \text{ of length } t \text{ that end in } s} P(Q)$$

Hidden Markov Models: Slide 13

Computation is exponential in t

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What is $P(q_t = s)$? Clever answer

• For each state s_i, define

$$p_t(i)$$
 = Prob. state is s_i at time t
= $P(q_t = s_i)$

· Easy to do inductive definition

$$\forall i \quad p_0(i) =$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

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What is $P(q_t = s)$? Clever answer

• For each state s_i , define $p_t(i)$ = Prob. state is s_i at time t

$$= P(q_t = s_i)$$

· Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

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Hidden Markov Models: Slide 15

What is $P(q_t = s)$? Clever answer

• For each state s_i, define

$$p_t(i)$$
 = Prob. state is s_i at time t
= $P(q_t = s_i)$

· Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

$$\sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) =$$

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What is $P(q_t = s)$? Clever answer

- For each state s_i, define $p_t(i)$ = Prob. state is s_i at time t $= P(q_t = s_i)$
- · Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) = \\ \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \\ \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \\ \sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \\ \sum_{i=1}^{N} a_{ij} p_t(i)$$

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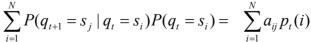
Hidden Markov Models: Slide 17

What is $P(q_t = s)$? Clever answer

- For each state s_i, define $p_t(i)$ = Prob. state is s_i at time t $= P(q_t = s_i)$
- · Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) = \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \sum_{i=1}^{N} P(q_{t+1} = s_i) = \sum_{i=1}^{N$$



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Hidden Markov Models: Slide 18

| order: | | | | | | | |
|--------|---|----------|----------|-----|----------|--|--|
| | | | | | | | |
| | t | $p_t(1)$ | $p_t(2)$ | / | $p_t(N)$ | | |
| | 0 | 0 — | 1 | 1/2 | 0 | | |
| | 1 | | | | | | |
| | : | 4 | | | | | |

Computation is simple.

What is $P(q_t = s)$? Clever answer

 $\sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \sum_{i=1}^{N} a_{ij} p_t(i)$

- For each state s_i , define $p_t(i)$ = Prob. state is s_i at time t = $P(q_t = s_i)$
- Easy to do inductive definition

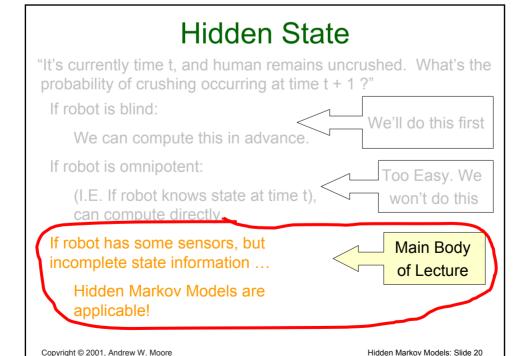
$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

$$\sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) =$$

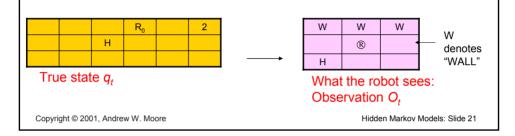
- Cost of computing P_t(i) for all states S_i is now O(t N²)
- The stupid way was O(N^t)
- This was a simple example
- It was meant to warm you up to this trick, called *Dynamic Programming*, because HMMs do many tricks like this.

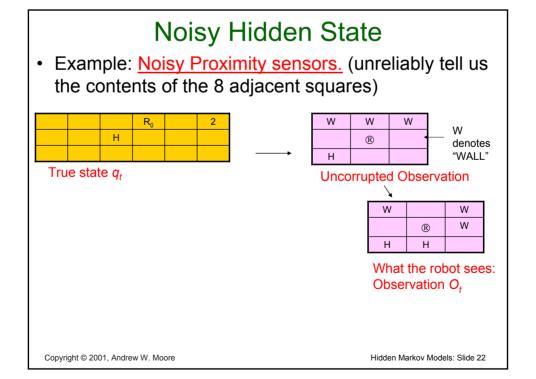
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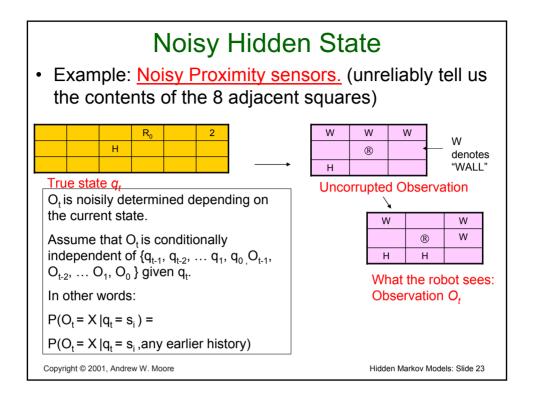


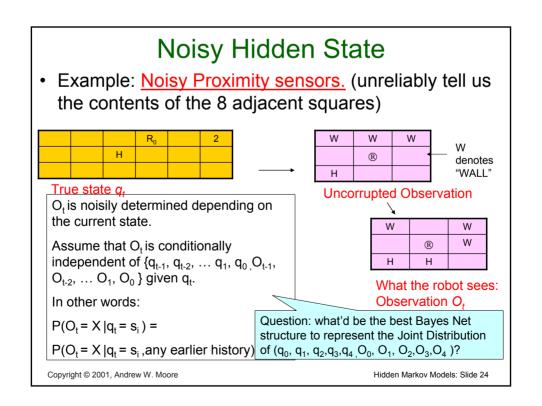
Hidden State

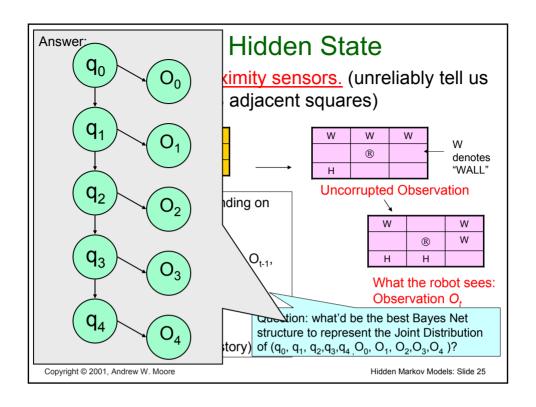
- The previous example tried to estimate $P(q_t = s_i)$ unconditionally (using no observed evidence).
- Suppose we can observe something that's affected by the true state.
- Example: <u>Proximity sensors.</u> (tell us the contents of the 8 adjacent squares)

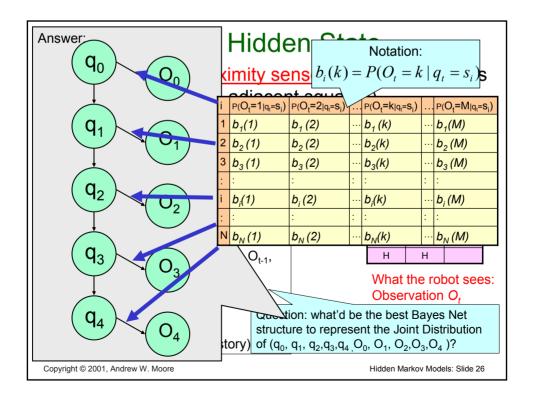












Hidden Markov Models

Our robot with noisy sensors is a good example of an HMM

Question 1: State Estimation

What is $P(q_T=S_i \mid O_1O_2...O_T)$

It will turn out that a new cute D.P. trick will get this for us.

· Question 2: Most Probable Path

Given $O_1O_2...O_T$, what is the most probable path that I took? And what is that probability?

Yet another famous D.P. trick, the VITERBI algorithm, gets this

Question 3: Learning HMMs:

Given $O_1O_2...O_T$, what is the maximum likelihood HMM that could have produced this string of observations?

Very very useful. Uses the E.M. Algorithm

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Hidden Markov Models: Slide 27

Are H.M.M.s Useful?

You bet !!

- Robot planning + sensing when there's uncertainty (e.g. Reid Simmons / Sebastian Thrun / Sven Koenig)
- Robot learning control (e.g. Yangsheng Xu's work)
- Speech Recognition/Understanding
 Phones → Words, Signal → phones
- Human Genome Project
 Complicated stuff your lecturer knows nothing about.
- · Consumer decision modeling
- Economics & Finance.

Plus at least 5 other things I haven't thought of.

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HMM Notation (from Rabiner's Survey)

The states are labeled S₁ S₂ .. S_N

For a particular trial....

be the number of observations Let T

is also the number of states passed through

 $O = O_1 O_2 ... O_T$ is the sequence of observations

 $Q = q_1 q_2 ... q_T$ is the notation for a path of states

 $\lambda = \langle N, M, \{\pi_i\}, \{a_{ii}\}, \{b_i(j)\} \rangle$ is the specification of an **HMM**

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Hidden Markov Models: Slide 29

HMM Formal Definition

 $b_N(M)$

An HMM, λ, is a 5-tuple consisting of

- N the number of states
- M the number of possible observations
- $\{\pi_1, \pi_2, ... \pi_N\}$ The starting state probabilities

 $P(q_0 = S_i) = \pi_i$ a_{22} a_{1N}

a₁₁ **a**₂₁ **a**₂₂ a_{2N} a_{NN}

 a_{N2}

b₁(2) ... $b_1(M)$ $b_1(1)$ $b_2(1)$ $b_2(2)$... $b_2(M)$

 $b_{N}(2)$

This is new. In our previous example, start state was deterministic

The state transition probabilities $P(q_{t+1}=S_i | q_t=S_i)=a_{ii}$

The observation probabilities $P(O_t=k \mid q_t=S_i)=b_i(k)$

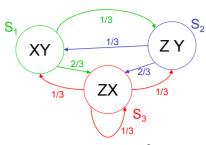
Hidden Markov Models: Slide 30

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 a_{N1}

 $b_N(1)$

Here's an HMM



Start randomly in state 1 or 2

Choose one of the output symbols in each state at random.

N = 3

M = 3

 $\pi_1 = 1/2$

 $\pi_2 = 1/2$

 $\pi_3 = 0$

 $a_{11} = 0$

 $a_{12} = 1/3$ $a_{22} = 0$ $a_{13} = 2/3$

 $a_{12} = 1/3$ $a_{13} = 1/3$

 $a_{22} = 0$ $a_{32} = 1/3$

 $a_{13} = 2/3$ $a_{13} = 1/3$

 $b_1(Y) = 1/2$

 $b_1(Z) = 0$

 $b_1(X) = 1/2$ $b_2(X) = 0$ $b_3(X) = 1/2$

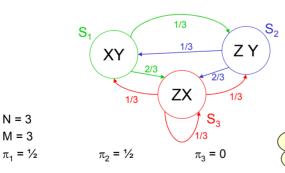
 $b_2(Y) = 1/2$ $b_3(Y) = 0$

 $b_2(Z) = 1/2$ $b_3(Z) = 1/2$

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Hidden Markov Models: Slide 31

Here's an HMM



Start randomly in state 1 or 2

Choose one of the output symbols in each state at random.

Let's generate a sequence of observations:

50-50 choice between S₁ and S₂

 $b_1(X) = \frac{1}{2}$ $b_1(Y) = \frac{1}{2}$

 $a_{11} = 0$ $a_{12} = \frac{1}{3}$

 $a_{13} = \frac{1}{3}$

 $b_2(X) = 0$

 $b_3(X) = \frac{1}{2}$

 $b_1(Z) = 0$ $b_2(Z) = \frac{1}{2}$

b₁ (Z) = 0

 $b_2(Z) = \frac{1}{2}$ $b_3(Z) = \frac{1}{2}$

 $a_{13} = \frac{2}{3}$

 $a_{13} = \frac{2}{3}$ $a_{13} = \frac{1}{3}$

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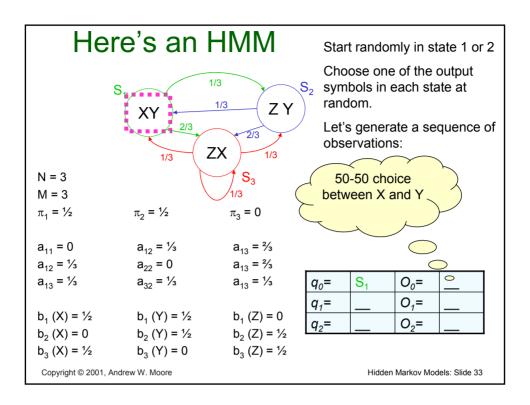
 $a_{12} = \frac{1}{3}$

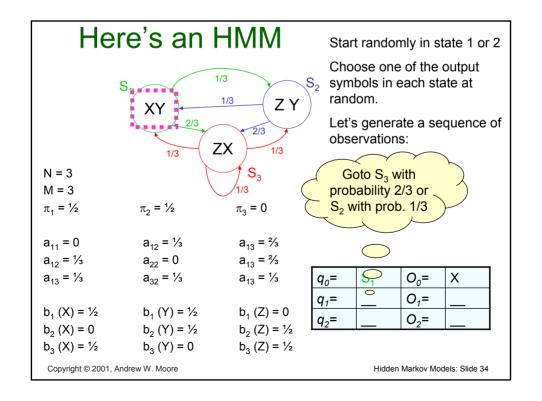
 $a_{22} = 0$

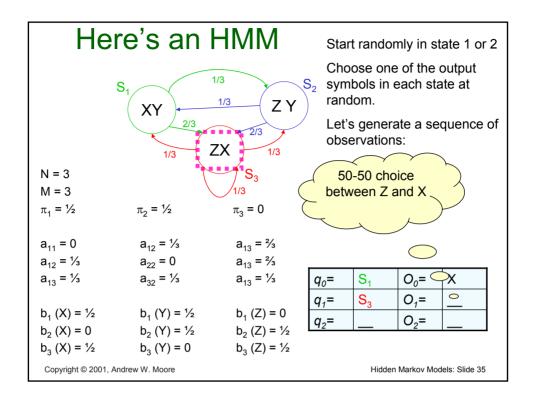
 $a_{32} = \frac{1}{3}$

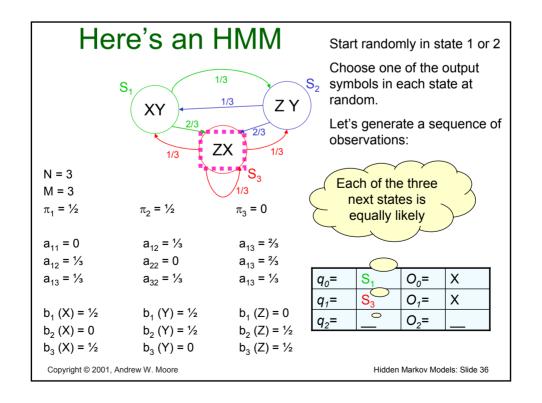
 $b_2(Y) = \frac{1}{2}$

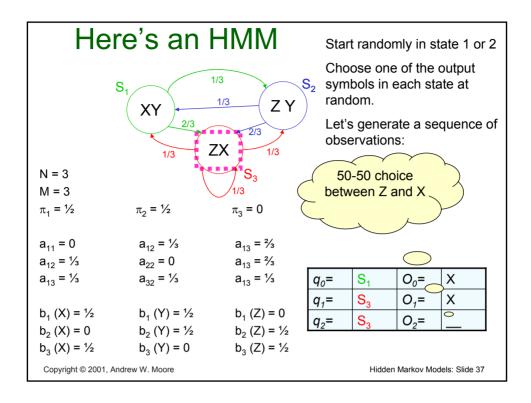
 $b_3(Y) = 0$

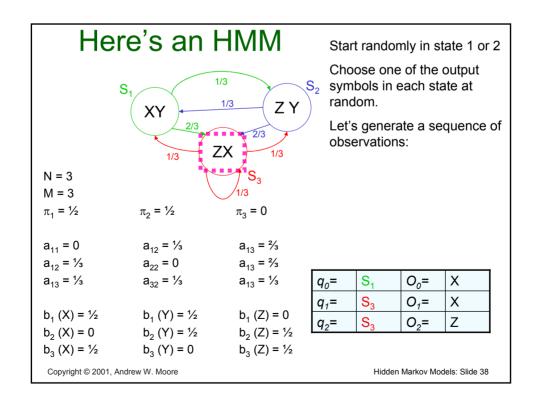




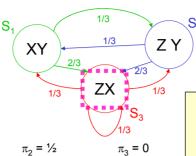












M = 3 $\pi_1 = \frac{1}{2}$ $\pi_2 = \frac{1}{2}$

N = 3

$$a_{11} = 0$$
 $a_{12} = \frac{1}{3}$ $a_{13} = \frac{2}{3}$ $a_{12} = \frac{1}{3}$ $a_{22} = 0$ $a_{13} = \frac{2}{3}$ $a_{13} = \frac{1}{3}$ $a_{32} = \frac{1}{3}$ $a_{13} = \frac{1}{3}$

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Start randomly in state 1 or 2

Choose one of the output symbols in each state at random.

Let's generate a sequence of observations:

This is what the observer has to work with...

| | 7 | | |
|------------------|---|------------------|---|
| q_0 = | ? | O ₀ = | Χ |
| $q_1 =$ | ? | O ₁ = | Χ |
| q ₂ = | ? | O ₂ = | Z |

Hidden Markov Models: Slide 39

Prob. of a series of observations

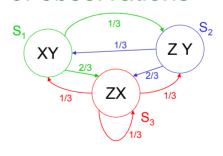
What is
$$P(\mathbf{O}) = P(O_1 O_2 O_3) = P(O_1 = X ^O_2 = X ^O_3 = Z)$$
?

Slow, stupid way:

$$P(\mathbf{O}) = \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \wedge \mathbf{Q})$$
$$= \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \mid \mathbf{Q}) P(\mathbf{Q})$$

How do we compute P(Q) for an arbitrary path Q?

How do we compute P(O|Q) for an arbitrary path Q?



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Prob. of a series of observations

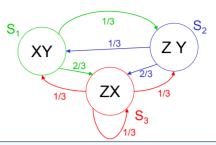
What is
$$P(\mathbf{O}) = P(O_1 O_2 O_3) = P(O_1 = X ^ O_2 = X ^ O_3 = Z)$$
?

Slow, stupid way:

$$P(\mathbf{O}) = \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \wedge \mathbf{Q})$$
$$= \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \mid \mathbf{Q}) P(\mathbf{Q})$$

How do we compute P(Q) for an arbitrary path Q?

How do we compute P(O|Q)for an arbitrary path Q?



 $P(Q) = P(q_1, q_2, q_3)$

 $=P(q_1) P(q_2,q_3|q_1)$ (chain rule)

 $=P(q_1) P(q_2|q_1) P(q_3|q_2,q_1)$ (chain)

 $=P(q_1) P(q_2|q_1) P(q_3|q_2)$ (why?)

Example in the case $Q = S_1 S_3 S_3$:

=1/2 * 2/3 * 1/3 = 1/9

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Hidden Markov Models: Slide 41

Prob. of a series of observations

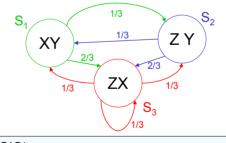
What is
$$P(\mathbf{O}) = P(O_1 O_2 O_3) = P(O_1 = X ^ O_2 = X ^ O_3 = Z)$$
?

Slow, stupid way:

$$P(\mathbf{O}) = \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \wedge \mathbf{Q})$$
$$= \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \mid \mathbf{Q}) P(\mathbf{Q})$$

How do we compute P(Q) for $= P(O_1 O_2 O_3 | q_1 q_2 q_3)$ an arbitrary path Q?

How do we compute P(O|Q) for an arbitrary path Q?



P(O|Q)

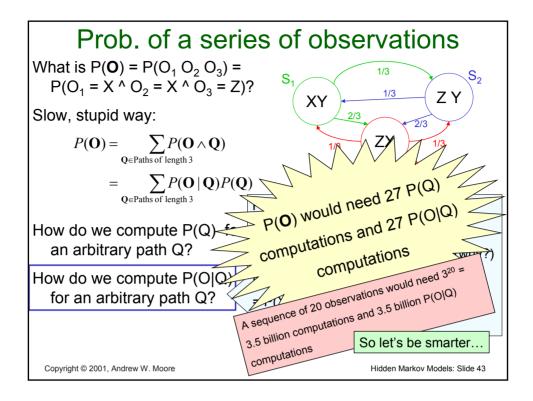
 $= P(O_1 | q_1) P(O_2 | q_2) P(O_3 | q_3) (why?)$

Example in the case $Q = S_1 S_3 S_3$:

 $= P(X|S_1) P(X|S_2) P(Z|S_3) =$

=1/2 * 1/2 * 1/2 = 1/8

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The Prob. of a given series of observations, non-exponential-cost-style

Given observations O₁ O₂ ... O_T

Define

$$\alpha_t(i) = P(O_1 O_2 ... O_t \land q_t = S_i \mid \lambda)$$
 where $1 \le t \le T$

 $\alpha_t(i)$ = Probability that, in a random trial,

- We'd have seen the first t observations
- ullet We'd have ended up in S_i as the t'th state visited.

In our example, what is $\alpha_2(3)$?

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$\alpha_t(i)$: easy to define recursively

 $\alpha_{t}(i) = P(O_{1} \ O_{2} \ \dots \ O_{T} \ \land \ q_{t} = S_{i} \mid \lambda) \ (\alpha_{t}(i) \ \text{can be defined stupidly by considering all paths length "t". How?})$

$$\alpha_{1}(i) = P(O_{1} \land q_{1} = S_{i})$$

$$= P(q_{1} = S_{i})P(O_{1}|q_{1} = S_{i})$$

$$= what?$$

$$\alpha_{t+1}(j) = P(O_{1}O_{2}...O_{t}O_{t+1} \land q_{t+1} = S_{j})$$

$$= \sum_{i=1}^{N} P(O_{1}O_{2}...O_{t} \land q_{t} = S_{i} \land O_{t+1} \land q_{t+1} = S_{j})$$

$$= \sum_{i=1}^{N} P(O_{t+1}, q_{t+1} = S_{j}|O_{1}O_{2}...O_{t} \land q_{t} = S_{i})P(O_{1}O_{2}...O_{t} \land q_{t} = S_{i})$$

$$= \sum_{i} P(O_{t+1}, q_{t+1} = S_{j}|q_{t} = S_{i})\alpha_{t}(i)$$

$$= \sum_{i} P(q_{t+1} = S_{j}|q_{t} = S_{i})P(O_{t+1}|q_{t+1} = S_{j})\alpha_{t}(i)$$

$$= \sum_{i} Q_{t+1}(O_{t+1})\alpha_{t}(i)$$

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in our example $\alpha_{t}(i) = P(O_{1}O_{2}..O_{t} \land q_{t} = S_{i}|\lambda)$ $\alpha_{1}(i) = b_{i}(O_{1})\pi_{i}$ $\alpha_{t+1}(j) = \sum a_{ij}b_{j}(O_{t+1})\alpha_{t}(i)$ XY 2/3 ZX 1/3 ZX 1/3 ZX 3/3

WE SAW
$$O_1 O_2 O_3 = X X Z$$

$$\alpha_1(1) = \frac{1}{4}$$
 $\alpha_1(2) = 0$
 $\alpha_1(3) = 0$
 $\alpha_2(1) = 0$
 $\alpha_2(2) = 0$
 $\alpha_2(3) = \frac{1}{12}$
 $\alpha_3(1) = 0$
 $\alpha_3(2) = \frac{1}{72}$
 $\alpha_3(3) = \frac{1}{72}$

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Easy Question

We can cheaply compute

$$\alpha_t(i)=P(O_1O_2...O_t \land q_t=S_i)$$

(How) can we cheaply compute

$$P(O_1O_2...O_t)$$
 ?

(How) can we cheaply compute

$$P(q_t=S_i|O_1O_2...O_t)$$

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Easy Question

We can cheaply compute

$$\alpha_t(i)=P(O_1O_2...O_t \land q_t=S_i)$$

(How) can we cheaply compute

$$P(O_1O_2...O_t) ? \sum_{i=1}^{N} \alpha_i(i)$$

(How) can we cheaply compute

$$P(q_t=S_i|O_1O_2...O_t)$$

$$\frac{\alpha_{t}(i)}{\sum_{j=1}^{N} \alpha_{t}(j)}$$

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Most probable path given observations

What's most probable path given $O_1O_2...O_T$, i.e.

What is
$$\underset{Q}{\operatorname{argmax}} P(Q|O_1O_2...O_T)$$
?

Slow, stupid answer:

$$\underset{Q}{\operatorname{argmax}} \ P(Q|O_1O_2...O_T)$$

= argmax
$$\frac{P(O_1O_2...O_T|Q)P(Q)}{P(O_1O_2...O_T)}$$

$$= \underset{Q}{\operatorname{argmax}} P(O_1 O_2 ... O_T | Q) P(Q)$$

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Efficient MPP computation

We're going to compute the following variables:

$$\delta_t(i) = \max_{\substack{q_1 q_2 ... q_{t-1}}} P(q_1 \ q_2 \, ... \ q_{t-1} \wedge q_t = S_i \wedge O_1 \, ... \, O_t)$$

= The Probability of the path of Length t-1 with the maximum chance of doing all these things:

...OCCURING

and

...ENDING UP IN STATE S_i

and

...PRODUCING OUTPUT O₁...O_t

DEFINE: $mpp_t(i) = that path$

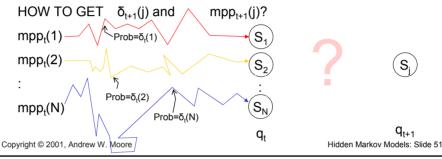
So: $\delta_t(i) = \text{Prob}(\text{mpp}_t(i))$

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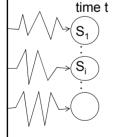
The Viterbi Algorithm

$$\begin{split} \delta_{t}(i) &= q_{1}q_{2}...q_{t-1} & \text{P}(q_{1}q_{2}...q_{t-1} \land q_{t} = S_{i} \land O_{1}O_{2}..O_{t}) \\ mpp_{t}(i) &= q_{1}q_{2}...q_{t-1} & \text{P}(q_{1}q_{2}...q_{t-1} \land q_{t} = S_{i} \land O_{1}O_{2}..O_{t}) \\ \delta_{1}(i) &= \text{one choice } \text{P}(q_{1} = S_{i} \land O_{1}) \\ &= \text{P}(q_{1} = S_{i}) \text{P}(O_{1}|q_{1} = S_{i}) \\ &= \pi.b.(O_{1}) \end{split}$$

Now, suppose we have all the $\delta_t(i)$'s and mpp $_t(i)$'s for all i.



The Viterbi Algorithm



time t+1

 S_i

The most prob path with last two states S_i S_i

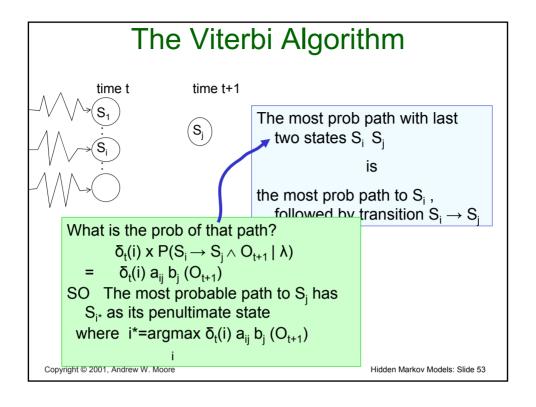
the most prob path to S_i, followed by transition $S_i \rightarrow S_i$

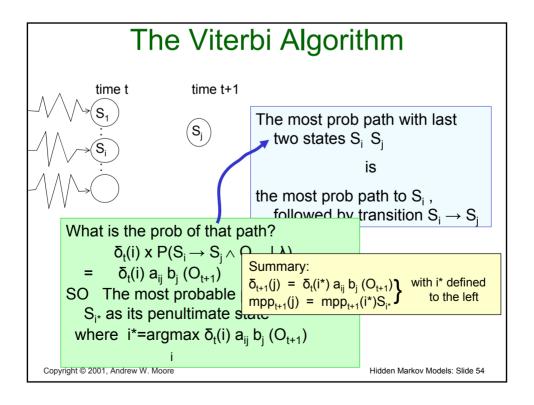
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 S_{i}

 q_{t+1}





What's Viterbi used for?

Classic Example

Speech recognition:

Signal → words

HMM → observable is signal

→ Hidden state is part of word formation

What is the most probable word given this signal?

UTTERLY GROSS SIMPLIFICATION

In practice: many levels of inference; not one big jump.

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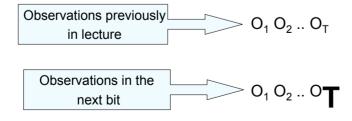
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HMMs are used and useful

But how do you design an HMM?

Occasionally, (e.g. in our robot example) it is reasonable to deduce the HMM from first principles.

But usually, especially in Speech or Genetics, it is better to infer it from large amounts of data. $O_1 O_2 ... O_T$ with a big "T".



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Inferring an HMM

Remember, we've been doing things like

$$P(O_1 O_2 .. O_T | \lambda)$$

That " λ " is the notation for our HMM parameters.

Now We have some observations and we want to estimate λ from them.

AS USUAL: We could use

- (i) MAX LIKELIHOOD $\lambda = \underset{\lambda}{\operatorname{argmax}} P(O_1 ... O_T | \lambda)$
- (ii) BAYES

Work out P($\lambda \mid O_1 ... O_T$)

and then take E[λ] or max P(λ | O₁ .. O_T)

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Max likelihood HMM estimation

Define

$$\begin{split} & \gamma_t(i) = P(q_t = S_i \mid O_1 O_2 ... O_T , \lambda) \\ & \epsilon_t(i,j) = P(q_t = S_i \land q_{t+1} = S_i \mid O_1 O_2 ... O_T , \lambda) \end{split}$$

 $\gamma_t(i)$ and $\epsilon_t(i,j)$ can be computed efficiently $\forall i,j,t$ (Details in Rabiner paper)

$$\sum_{t=1}^{T-1} \gamma_t(i) = \sum_{t=1}^{T-1} \gamma_t(i)$$
 Expected number of transitions out of state i during the path

$$\sum_{t=1}^{T-1} \mathcal{E}_t (i,j) = \text{ Expected number of transitions from state i to state j during the path}$$

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$$\begin{split} & \gamma_t(i) = \mathrm{P}\big(q_t = S_i \big| O_1 O_2 ... O_T, \lambda\big) \\ & \varepsilon_t(i,j) = \mathrm{P}\big(q_t = S_i \land q_{t+1} = S_j \big| O_1 O_2 ... O_T, \lambda\big) \\ & \sum_{t=1}^{T-1} \gamma_t(i) = \mathrm{expected\ number\ of\ transitions\ out\ of\ state\ i\ during\ path} \\ & \sum_{t=1}^{T-1} \varepsilon_t(i,j) = \mathrm{expected\ number\ of\ transitions\ out\ of\ i\ and\ into\ j\ during\ path} \end{split}$$

HMM estimation

Notice
$$\frac{\sum_{t=1}^{T-1} \varepsilon_t(i,j)}{\sum_{t=1}^{T-1} \gamma_t(i)} = \frac{\left(\begin{array}{c} \text{expected frequency} \\ i \to j \end{array}\right)}{\left(\begin{array}{c} \text{expected frequency} \\ i \end{array}\right)}$$

= Estimate of Prob(Next state S_i) This state S_i)

We can re-estimate

$$\mathbf{a}_{ij} \leftarrow \frac{\sum \varepsilon_t(i,j)}{\sum \gamma_t(i)}$$

We can also re - estimate

$$b_i(O_k) \leftarrow \cdots$$

(See Rabiner)

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EM for HMMs

If we knew λ we could estimate EXPECTATIONS of quantities such as

Expected number of times in state i

Expected number of transitions $i \rightarrow j$

If we knew the quantities such as

Expected number of times in state i

Expected number of transitions $i \rightarrow j$

We could compute the MAX LIKELIHOOD estimate of

$$\lambda = \langle \{a_{ij}\}, \{b_i(j)\}, \pi_i \rangle$$

Roll on the EM Algorithm...

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EM 4 HMMs

- 1. Get your observations O₁ ...O_T
- 2. Guess your first λ estimate $\lambda(0)$, t=0
- 3. t = t+1
- 4. Given $O_1 ... O_T$, $\lambda(t)$ compute $\gamma_t(i), \ \epsilon_t(i,j) \quad \forall \ 1 \le t \le T, \quad \forall \ 1 \le i \le N, \quad \forall \ 1 \le j \le N$
- 5. Compute expected freq. of state i, and expected freq. i→j
- 6. Compute new estimates of a_{ij} , $b_j(k)$, π_i accordingly. Call them $\lambda(t+1)$
- 7. Goto 3, unless converged.
- Also known (for the HMM case) as the BAUM-WELCH algorithm.

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Bad News

There are lots of local minima

Good News

 The local minima are usually adequate models of the data.

Notice

- EM does not estimate the number of states. That must be given.
- Often, HMMs are forced to have some links with zero probability. This is done by setting a_{ij} =0 in initial estimate $\lambda(0)$
- Easy extension of everything seen today: HMMs with real valued outputs

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What You Should Know

- · What is an HMM?
- Computing (and defining) α_t(i)
- · The Viterbi algorithm
- Outline of the EM algorithm
- To be very happy with the kind of maths and analysis needed for HMMs
- Fairly thorough reading of Rabiner* up to page 266*
 [Up to but not including "IV. Types of HMMs"].
- *L. R. Rabiner, "A Tutorial on Hidden Markov Models and Selected Applications in Speech Recognition," Proc. of the IEEE, Vol.77, No.2, pp.257--286, 1989.

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DON'T PANIC: starts on p. 257.